Parallel computing An introduction

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Outline

- Hardware architecture
- serial vs. parallel computing
- OpenMP parallelisation
- MPI parallelisation
- mixed OpenMP and MPI
- perfectly parallelizing algorithms
- hydrodynamics and domain decomposition
- long-range forces and communication

CPU and GPU



CPU with multiple GPUs



CPU network



Why parallel computing

- multiple cores are faster
 - if algorithm parallelises well
 - if communication is fast
- data does not fit on memory
- different tasks on different hardware
 - complex instructions on CPU
 - simple code on GPU

- shared memory parallelization
- one global process on the machine
- temporally occupies multiple cores/threads
- pro: simpler coding
- cons: limited to one node / CPU unit





time



- *distributed* memory parallelization
- multiple individual OS processes
- in priniple independent execution
- occupies multiple cores
- every process own part of memory
- processes need to communicate
- more complicated coding
- no limitations on *local* memory and *local* number of cores





MPI barrier: wait until all processes are here

basic structure of MPI program

#include <mpi.h>
#include <stdio.h>
int main(int argc, char** argv) {

// Initialize the MPI environment
MPI_Init(NULL, NULL);

// Get the number of processes
int world_size;
MPI_Comm_size(MPI_COMM_WORLD, &world_size);

// Get the rank of the process
int world_rank;
MPI_Comm_rank(MPI_COMM_WORLD, &world_rank);

```
// Finalize the MPI environment.
MPI_Finalize();
```

MPI communication

int main(int argc, char** argv) {

• • •

. . .

}

// send data to other processor
MPI_Send(data, count, datatype, destination, tag, MPI_communicator);

// receive data from other processor
MPI_Recv(data, count, datatype, source, tag, MPI_communicator, status);

MPI communication II

```
int main(int argc, char** argv) {
    ...
    // initialise 2 MPI processes
    if(world_rank == 0)
    {
        // send data to other processor
        MPI_Send(data, count, datatype, destination=1, tag, MPI_comm);
    }
    else // world_rank == 1
    {
        // receive data from other processor
        MPI_Recv(data, count, datatype, source=0, tag, MPI_comm, status);
    }
}
```

. . .

}

MPI communication III

```
int main(int argc, char** argv) {
    . . .
    // initialise N MPI processes
    if(world_rank == 0) {
        for(int i=1; i<N; i++)</pre>
        {
            // send data to other processor
            MPI_Send(data, count, datatype, destination=i, tag, MPI_comm);
        }
    }
    else
    {
        // receive data from other processor
        MPI_Recv(data, count, datatype, source=0, tag, MPI_comm, status);
    }
```

. . .

}

MPI communication IV

```
int main(int argc, char** argv) {
    ...
    // distribute to all
    if(world_rank == 0) {
        // send data to other processor
        MPI_Bcast(data, count, datatype, tag, MPI_comm);
    }
    else
    {
        // receive data from other processor
        MPI_Recv(data, count, datatype, source=0, tag, MPI_comm, status);
    }
```

. . .

}

MPI other commands

```
int main(int argc, char** argv) {
    . . .
    // wait here
    MPI_Barrier(...)
    // collect from all
    MPI_Gather(...)
    // reduce
    MPI_Reduce(..., mode=MODE)
    . . .
}
MODE:
MPI_MAX : find minimum
MPI_MIN : find maximum
MPI_SUM : sum all values
MPI_PROD : multiply all values
MPI_LAND : logical and
MPI_LOR : logical or
. . . .
```



option 1

- simple pinning
- each process: one core



option 2

- simple pinning
- processes distributed onto several nodes
- 2/6 cores occupied
- each process 1/2 memory

mixed OpenMP and MPI

- start MPI process on every node
- inside node OpenMP with shared memory

Perfect parallelization

- vector multiplication
- simple matrix operations
- Monte Carlo simulations
- independent parameter scan

hydrodynamics everywhere







hydrodynamics

- Solve discretized fluid equations on a grid
- simplest case: uniform, periodic grid



hydrodynamics

- Solve discretized fluid equations on a grid
- simplest case: uniform, periodic grid
- split domain between processors



hydrodynamcis - equations



discretization

temporal and spatial discretization

$$\frac{\partial y}{\partial t} \rightarrow \frac{y_i^{n+1} - y_i^n}{\Delta t}$$
$$\frac{\partial^2 y}{\partial x^2} \rightarrow \frac{y_{i+1}^n - 2y_i^n + y_{i-1}^n}{\Delta x^2}$$

 3-point stencil: need one neighbour in each direction

need neighbour cells





how to communicate

stupid way

do time_loop

```
do x_loop
  do y_loop
    if at boundary
        # communicate
        MPI_get_neighbour()
        density = ...
        momentum = ...
        energy = ...
        done
        done
        done
```

clever way

do time_loop
 # get a copy of neighbours
 MPI_get_neighbours()
 do x_loop
 do y_loop

```
density = ...
momentum = ...
energy = ...
done
done
MPI_send_neighbours()
done
```

guard (ghost) cells



use guard cells to reduce communication

galaxy and ISM



galaxy and ISM



Lifecycle of molecular clouds

Cooling & Collapse



Stellar Feedback & Outflows

simulations of interstellar gas

- dense cold gas that forms clouds and stars
- diffuse warm gas
- hot gas that escapes the galaxy
- different scales (space/time), so need adaptive grid



adaptive mesh refinement



example: Sedov explosion



- dynamically follow the interesting gas structures (here shock)
 refine and derefine
- more complicated domain decomposition
- dynamical redistribution of regions between the cores

Domain decomposition

- simplest way:
 - each processor same number of cells
 - select domain with least communication (shortest border)
 - perfect memory distribution
- problematic if different cells require different cost
 - iterations depend on density, temperature
 - iterations depend on position

memory balancing



local time steps (actual work)





physics load balancing

- perfect distribution in memory
- but small cells interact on smaller time scales
- small blocks need to do more iterations
- cores with small cells do more work!

$dt = dt_0$	dt ₀ /2	dt ₀ /4
	dt ₀ /2	dt ₀ /2
dt = dt ₀	dt = 6	dt _o

another example: ISM

density

local hydro time



time scales

density

free-fall time



NN vs. long-range interaction

hydro (NN) cell-by-cell speed









gravity (long-range) instantaneous speed



direct neighbour vs. long-range

hydro (dir. neighbour) cell-by-cell speed

- one guard cell works
- two in case of 5pt stencil
- small additional memory
- communication to neighbour processors, globally asynchronous

gravity (long-range) instantaneous speed

- every cell depends on every cell (N²)
- every processor need entire grid information
- reduce information, approximate computation
- tree methods, particlemesh methods
- still communication accross all processors

simple example: tree gravity

- reduce objects at large distances to centre of mass
- compute force
 between centres



simple example: tree gravity

- reduce objects at large distances to centre of mass
- compute force
 between centres
- close clouds need direct integration





if tree structure is known:

- pro: efficient communication with necessary processors
- con: every process needs to have tree information
 - --> tree needs to be communicated

scaling

- strong scaling: how the solution time varies with the number of processors for a fixed total problem size
 - -ideal: t = N_{tot}/N_{proc} , speedup = 1/t = N_{proc}/N_{tot}
- weak scaling: how the solution time varies with the number of processors for a *fixed problem size per processor* – ideal: t = const, speedup = const

scaling in real application



ISM simulations



ISM simulations



processes

- MHD (local)
- self-gravity* (tree)
- external potential (analytic)
- radiation* and shielding* (tree)
- in practice:
 - tree efficient in terms of comp. cost
 - tree stores variables for (*), a lot of memory
 - sim. "memory limited"
 - more cores *would* help, but not enough memory

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- reading data:
 - few thousand cores direct reading efficient caches -> OK
 - one process reads -> MPI distribution
- writing data:
 - parallel writing at random positions in file: data race! (only one process allowed, lock)
 - files split like domain decomposition (every processor separate file with local data)
 one process: MPI collection -> writing

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code

- FLASH / Arepo
- C / C++ / Fortran
- MPI / MPI+OpenMP
- ca. 400.000 lines
- problem: 100.000 lines
- current sim: 40 Mio CPUh, 250 TB

problems and conclusions

- all computations must be parallel
- many runs need >1000 cores
- MPI and combined MPI/openMP
- most of work:
 - numerical methods for the physics equations
 - optimization and efficient parallelization
- so far missing: machine learning methods